

Chapter 3

Sorting:

Sorting refers to arranging data in a particular format. Sorting algorithm specifies the way to arrange data in a particular order. Most common orders are in numerical or lexicographical order. The importance of sorting lies in the fact that data searching can be optimized to a very high level, if data is stored in a sorted manner. Sorting is also used to represent data in more readable formats. Following are some of the examples of sorting in real-life scenarios –

Telephone Directory – The telephone directory stores the telephone numbers of people sorted by their names, so that the names can be searched easily.

Dictionary – The dictionary stores words in an alphabetical order so that searching of any word becomes easy.

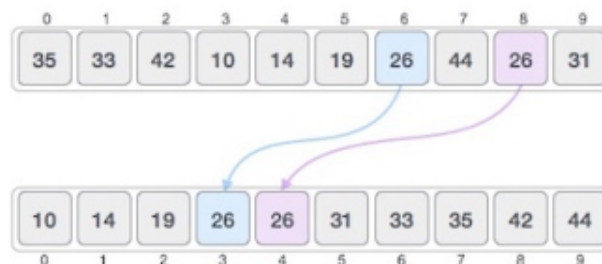
In-place Sorting and Not-in-place Sorting:

Sorting algorithms may require some extra space for comparison and temporary storage of few data elements. In-place sorting algorithms do not require any extra space and sorting is said to happen in-place within the array itself. Bubble sort is an example of in-place sorting.

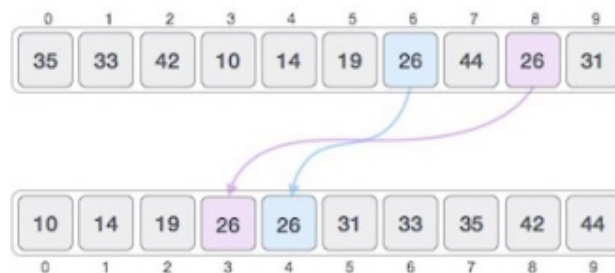
However, in some sorting algorithms, the program requires space which is more than or equal to the elements being sorted. Sorting which uses equal or more space is called not-in-place sorting. Merge-sort is an example of not-in-place sorting.

Stable and Unstable Sorting:

If a sorting algorithm, after sorting the contents, does not change the sequence of similar content in which they appear, it is called stable sorting.



If a sorting algorithm, after sorting the contents, changes the sequence of similar content in which they appear, it is called unstable sorting.



Stability of an algorithm matters when we wish to maintain the sequence of original elements, like in a tuple for example.

Adaptive and Non-Adaptive Sorting Algorithm:

A sorting algorithm is said to be adaptive, if it takes advantage of already 'sorted' elements in the list that is to be sorted. That is, while sorting if the source list has some element already sorted, adaptive algorithms will take this into account and will try not to re-order them.

A non-adaptive algorithm is one which does not take into account the elements which are already sorted. They try to force every single element to be re-ordered to confirm their sortedness.

Important Terms

Some terms are generally coined while discussing sorting techniques, here is a brief introduction to them –

Increasing Order:

A sequence of values is said to be in increasing order, if the successive element is greater than the previous one. For example, 1, 3, 4, 6, 8, 9 are in increasing order, as every next element is greater than the previous element.

Decreasing Order:

A sequence of values is said to be in decreasing order, if the successive element is less than the current one. For example, 9, 8, 6, 4, 3, 1 are in decreasing order, as every next element is less than the previous element.

Non-Increasing Order:

A sequence of values is said to be in non-increasing order, if the successive element is less than or equal to its previous element in the sequence. This order occurs when the sequence contains duplicate values. For example, 9, 8, 6, 3, 3, 1 are in non-increasing order, as every next element is less than or equal to (in case of 3) but not greater than any previous element.

Non-Decreasing Order:

A sequence of values is said to be in non-decreasing order, if the successive element is greater than or equal to its previous element in the sequence. This order occurs when the sequence contains duplicate values. For example, 1, 3, 3, 6, 8, 9 are in non-decreasing order, as every next element is greater than or equal to (in case of 3) but not less than the previous one.

Bubble Sort:

Bubble sort is a simple sorting algorithm. This sorting algorithm is comparison-based algorithm in which each pair of adjacent elements is compared and the elements are swapped if they are not in order. This algorithm is not suitable for large data sets as its average and worst case complexity are of $O(n^2)$ where n is the number of items.

How Bubble Sort Works?

We take an unsorted array for our example. Bubble sort takes $O(n^2)$ time so we're keeping it short and precise.



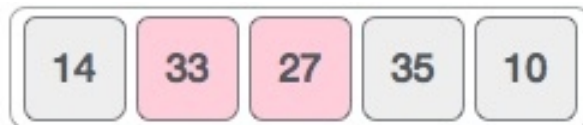
Bubble sort starts with very first two elements, comparing them to check which one is greater.



In this case, value 33 is greater than 14, so it is already in sorted locations. Next, we compare 33 with 27.



We find that 27 is smaller than 33 and these two values must be swapped.



The new array should look like this –



Next we compare 33 and 35. We find that both are in already sorted positions.



Then we move to the next two values, 35 and 10.



We know then that 10 is smaller 35. Hence they are not sorted.



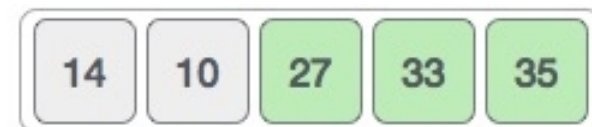
We swap these values. We find that we have reached the end of the array. After one iteration, the array should look like this –



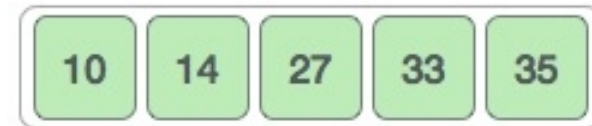
To be precise, we are now showing how an array should look like after each iteration. After the second iteration, it should look like this –



Notice that after each iteration, at least one value moves at the end.



And when there's no swap required, bubble sort learns that an array is completely sorted.



Algorithm:

We assume list is an array of n elements. We further assume that swap function swaps the values of the given array elements.

```
begin BubbleSort(list)
```

```
  for all elements of list
    if list[i] > list[i+1]
      swap(list[i], list[i+1])
    end if
  end for
```

```
  return list
```

```
end BubbleSort
```

C program for Bubble Sort:

```
#include <stdio.h>
```

```
#include <stdbool.h>
```

```
#define MAX 10
```

```
int list[MAX] = {1,8,4,6,0,3,5,2,7,9};
```

```

void display() {
    int i;
    printf("[");

    // navigate through all items
    for(i = 0; i < MAX; i++) {
        printf("%d ",list[i]);
    }

    printf("]\n");
}

void bubbleSort() {
    int temp;
    int i,j;

    bool swapped = false;

    // loop through all numbers
    for(i = 0; i < MAX-1; i++) {
        swapped = false;

        // loop through numbers falling ahead
        for(j = 0; j < MAX-1-i; j++) {
            printf("    Items compared: [ %d, %d ] ", list[j],list[j+1]);

            // check if next number is lesser than current no
            // swap the numbers.
            // (Bubble up the highest number)

            if(list[j] > list[j+1]) {
                temp = list[j];
                list[j] = list[j+1];
                list[j+1] = temp;

                swapped = true;
                printf("=> swapped [%d, %d]\n",list[j],list[j+1]);
            }else {
                printf("=> not swapped\n");
            }
        }

        // if no number was swapped that means
        // array is sorted now, break the loop.
        if(!swapped) {
            break;
        }
    }
}

```

```

    }

    printf("Iteration %d#: ",(i+1));
    display();
}

}

main() {
    printf("Input Array: ");
    display();
    printf("\n");

    bubbleSort();
    printf("\nOutput Array: ");
    display();
}

```

If we compile and run the above program, it will produce the following result –

Input Array: [1 8 4 6 0 3 5 2 7 9]

```

    Items compared: [ 1, 8 ] => not swapped
    Items compared: [ 8, 4 ] => swapped [4, 8]
    Items compared: [ 8, 6 ] => swapped [6, 8]
    Items compared: [ 8, 0 ] => swapped [0, 8]
    Items compared: [ 8, 3 ] => swapped [3, 8]
    Items compared: [ 8, 5 ] => swapped [5, 8]
    Items compared: [ 8, 2 ] => swapped [2, 8]
    Items compared: [ 8, 7 ] => swapped [7, 8]
    Items compared: [ 8, 9 ] => not swapped
Iteration 1#: [1 4 6 0 3 5 2 7 8 9 ]
    Items compared: [ 1, 4 ] => not swapped
    Items compared: [ 4, 6 ] => not swapped
    Items compared: [ 6, 0 ] => swapped [0, 6]
    Items compared: [ 6, 3 ] => swapped [3, 6]
    Items compared: [ 6, 5 ] => swapped [5, 6]
    Items compared: [ 6, 2 ] => swapped [2, 6]
    Items compared: [ 6, 7 ] => not swapped
    Items compared: [ 7, 8 ] => not swapped
Iteration 2#: [1 4 0 3 5 2 6 7 8 9 ]
    Items compared: [ 1, 4 ] => not swapped
    Items compared: [ 4, 0 ] => swapped [0, 4]
    Items compared: [ 4, 3 ] => swapped [3, 4]
    Items compared: [ 4, 5 ] => not swapped
    Items compared: [ 5, 2 ] => swapped [2, 5]
    Items compared: [ 5, 6 ] => not swapped
    Items compared: [ 6, 7 ] => not swapped

```

Iteration 3#: [1 0 3 4 2 5 6 7 8 9]
 Items compared: [1, 0] => swapped [0, 1]
 Items compared: [1, 3] => not swapped
 Items compared: [3, 4] => not swapped
 Items compared: [4, 2] => swapped [2, 4]
 Items compared: [4, 5] => not swapped
 Items compared: [5, 6] => not swapped
 Iteration 4#: [0 1 3 2 4 5 6 7 8 9]
 Items compared: [0, 1] => not swapped
 Items compared: [1, 3] => not swapped
 Items compared: [3, 2] => swapped [2, 3]
 Items compared: [3, 4] => not swapped
 Items compared: [4, 5] => not swapped
 Iteration 5#: [0 1 2 3 4 5 6 7 8 9]
 Items compared: [0, 1] => not swapped
 Items compared: [1, 2] => not swapped
 Items compared: [2, 3] => not swapped
 Items compared: [3, 4] => not swapped

Output Array: [0 1 2 3 4 5 6 7 8 9]

Selection Sort: Selection sort is a simple sorting algorithm. This sorting algorithm is an in-place comparison-based algorithm in which the list is divided into two parts, the sorted part at the left end and the unsorted part at the right end. Initially, the sorted part is empty and the unsorted part is the entire list. The smallest element is selected from the unsorted array and swapped with the leftmost element, and that element becomes a part of the sorted array. This process continues moving unsorted array boundary by one element to the right. This algorithm is not suitable for large data sets as its average and worst case complexities are of $O(n^2)$, where n is the number of items.

How Selection Sort Works?

Consider the following depicted array as an example.



For the first position in the sorted list, the whole list is scanned sequentially. The first position where 14 is stored presently, we search the whole list and find that 10 is the lowest value.



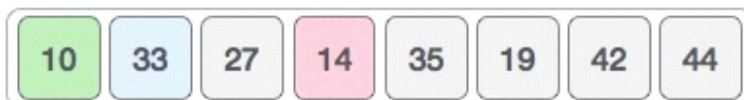
So we replace 14 with 10. After one iteration 10, which happens to be the minimum value in the list, appears in the first position of the sorted list.



For the second position, where 33 is residing, we start scanning the rest of the list in a linear manner.



We find that 14 is the second lowest value in the list and it should appear at the second place. We swap these values.



After two iterations, two least values are positioned at the beginning in a sorted manner.



The same process is applied to the rest of the items in the array.

Following is a pictorial depiction of the entire sorting process –



Algorithm:

Step 1 – Set MIN to location 0
Step 2 – Search the minimum element in the list
Step 3 – Swap with value at location MIN
Step 4 – Increment MIN to point to next element
Step 5 – Repeat until list is sorted

C program for Selection Sort:

```
#include <stdio.h>
#include <stdbool.h>
#define MAX 7

int intArray[MAX] = {4,6,3,2,1,9,7};

void printline(int count) {
    int i;

    for(i = 0; i < count-1; i++) {
        printf("=");
    }

    printf("\n");
}

void display() {
    int i;
    printf("[");

    // navigate through all items
    for(i = 0; i < MAX; i++) {
        printf("%d ", intArray[i]);
    }

    printf("]\n");
}

void selectionSort() {
    int indexMin, i, j;

    // loop through all numbers
    for(i = 0; i < MAX-1; i++) {

        // set current element as minimum
        indexMin = i;

        // check the element to be minimum
        for(j = i+1; j < MAX; j++) {
```

```

        if(intArray[j] < intArray[indexMin]) {
            indexMin = j;
        }
    }

    if(indexMin != i) {
        printf("Items swapped: [ %d, %d ]\n" , intArray[i], intArray[indexMin]);

        // swap the numbers
        int temp = intArray[indexMin];
        intArray[indexMin] = intArray[i];
        intArray[i] = temp;
    }

    printf("Iteration %d#:",(i+1));
    display();
}

main() {
    printf("Input Array: ");
    display();
    printline(50);
    selectionSort();
    printf("Output Array: ");
    display();
    printline(50);
}

```

If we compile and run the above program, it will produce the following result –
 Input Array: [4 6 3 2 1 9 7]

```

=====
Items swapped: [ 4, 1 ]
Iteration 1#:[1 6 3 2 4 9 7 ]
Items swapped: [ 6, 2 ]
Iteration 2#:[1 2 3 6 4 9 7 ]
Iteration 3#:[1 2 3 6 4 9 7 ]
Items swapped: [ 6, 4 ]
Iteration 4#:[1 2 3 4 6 9 7 ]
Iteration 5#:[1 2 3 4 6 9 7 ]
Items swapped: [ 9, 7 ]
Iteration 6#:[1 2 3 4 6 7 9 ]
Output Array: [1 2 3 4 6 7 9 ]

```

Merge Sort: Merge sort is a sorting technique based on divide and conquer technique. With worst-case time complexity being $O(n \log n)$, it is one of the most respected algorithms. Merge sort first divides the array into equal halves and then combines them in a sorted manner.

How Merge Sort Works?

To understand merge sort, we take an unsorted array as the following –



We know that merge sort first divides the whole array iteratively into equal halves unless the atomic values are achieved. We see here that an array of 8 items is divided into two arrays of size 4.



This does not change the sequence of appearance of items in the original. Now we divide these two arrays into halves.



We further divide these arrays and we achieve atomic value which can no more be divided



Atul starts from here

Now, we combine them in exactly the same manner as they were broken down. Please note the color codes given to these lists.

We first compare the element for each list and then combine them into another list in a sorted manner. We see that 14 and 33 are in sorted positions. We compare 27 and 10 and in the target list of 2 values we put 10 first, followed by 27. We change the order of 19 and 35 whereas 42 and 44 are placed sequentially.



In the next iteration of the combining phase, we compare lists of two data values, and merge them into a list of found data values placing all in a sorted order.



After the final merging, the list should look like this –



Now we should learn some programming aspects of merge sorting.

Algorithm:

Merge sort keeps on dividing the list into equal halves until it can no more be divided. By definition, if it is only one element in the list, it is sorted. Then, merge sort combines the smaller sorted lists keeping the new list sorted too.

Step 1 – if it is only one element in the list it is already sorted, return.

Step 2 – divide the list recursively into two halves until it can no more be divided.

Step 3 – merge the smaller lists into new list in sorted order.

C Program for Merge Sort:

```
#include <stdio.h>
```

```
#define max 10
```

```
int a[10] = {10, 14, 19, 26, 27, 31, 33, 35, 42, 44 };
```

```
int b[10];
```

```
void merging(int low, int mid, int high) {
```

```
    int l1, l2, i;
```

```
    for(l1 = low, l2 = mid + 1, i = low; l1 <= mid && l2 <= high; i++) {
```

```
        if(a[l1] <= a[l2])
```

```
            b[i] = a[l1++];
```

```
        else
```

```
            b[i] = a[l2++];
```

```
    }
```

```
    while(l1 <= mid)
```

```
        b[i++] = a[l1++];
```

```
    while(l2 <= high)
```

```
        b[i++] = a[l2++];
```

```
    for(i = low; i <= high; i++)
```

```
        a[i] = b[i];
```

```
}
```

```
void sort(int low, int high) {
```

```
    int mid;
```

```
    if(low < high) {
```

```

        mid = (low + high) / 2;
        sort(low, mid);
        sort(mid+1, high);
        merging(low, mid, high);
    }else {
        return;
    }
}

int main() {
    int i;

    printf("List before sorting\n");

    for(i = 0; i <= max; i++)
        printf("%d ", a[i]);

    sort(0, max);

    printf("\nList after sorting\n");

    for(i = 0; i <= max; i++)
        printf("%d ", a[i]);
}

```

If we compile and run the above program, it will produce the following result –

```

Output
List before sorting
10 14 19 26 27 31 33 35 42 44 0
List after sorting
0 10 14 19 26 27 31 33 35 42 44

```

Insertion Sort: This is an in-place comparison-based sorting algorithm. Here, a sub-list is maintained which is always sorted. For example, the lower part of an array is maintained to be sorted. An element which is to be inserted in this sorted sub-list, has to find its appropriate place and then it has to be inserted there. Hence the name, insertion sort.

The array is searched sequentially and unsorted items are moved and inserted into the sorted sub-list (in the same array). This algorithm is not suitable for large data sets as its average and worst case complexity are of $O(n^2)$, where n is the number of items.

How Insertion Sort Works?

We take an unsorted array for our example.



Insertion sort compares the first two elements.



It finds that both 14 and 33 are already in ascending order. For now, 14 is in sorted sub-list.



Insertion sort moves ahead and compares 33 with 27.



And finds that 33 is not in the correct position.



It swaps 33 with 27. It also checks with all the elements of sorted sub-list. Here we see that the sorted sub-list has only one element 14, and 27 is greater than 14. Hence, the sorted sub-list remains sorted after swapping.



By now we have 14 and 27 in the sorted sub-list. Next, it compares 33 with 10.



These values are not in a sorted order.



So we swap them.



However, swapping makes 27 and 10 unsorted.



Hence, we swap them too.



Again we find 14 and 10 in an unsorted order.



We swap them again. By the end of third iteration, we have a sorted sub-list of 4 items.



This process goes on until all the unsorted values are covered in a sorted sub-list.

Algorithm:

Now we have a bigger picture of how this sorting technique works, so we can derive simple steps by which we can achieve insertion sort.

Step 1 – If it is the first element, it is already sorted. return 1

Step 2 – Pick next element

Step 3 – Compare with all elements in the sorted sub-list

Step 4 – Shift all the elements in the sorted sub-list that is greater than the value to be sorted

Step 5 – Insert the value

Step 6 – Repeat until list is sorted

C Program for Insertion Sort:

```
#include <stdio.h>
```

```
#include <stdbool.h>
```

```
#define MAX 7
```

```
int intArray[MAX] = {4,6,3,2,1,9,7};
```

```
void printline(int count) {  
    int i;
```

```

    for(i = 0;i <count-1;i++) {
        printf("=");
    }

    printf("\n");
}

void display() {
    int i;
    printf("[");

    // navigate through all items
    for(i = 0;i<MAX;i++) {
        printf("%d ",intArray[i]);
    }

    printf("]\n");
}

void insertionSort() {

    int valueToInsert;
    int holePosition;
    int i;

    // loop through all numbers
    for(i = 1; i < MAX; i++) {

        // select a value to be inserted.
        valueToInsert = intArray[i];

        // select the hole position where number is to be inserted
        holePosition = i;

        // check if previous no. is larger than value to be inserted
        while (holePosition > 0 && intArray[holePosition-1] > valueToInsert) {
            intArray[holePosition] = intArray[holePosition-1];
            holePosition--;
            printf(" item moved : %d\n" , intArray[holePosition]);
        }

        if(holePosition != i) {
            printf(" item inserted : %d, at position : %d\n" , valueToInsert,holePosition);
            // insert the number at hole position
            intArray[holePosition] = valueToInsert;
        }

        printf("Iteration %d#:",i);
    }
}

```

```

        display();
    }
}

main() {
    printf("Input Array: ");
    display();
    printline(50);
    insertionSort();
    printf("Output Array: ");
    display();
    printline(50);
}

```

If we compile and run the above program, it will produce the following result –

Input Array: [4 6 3 2 1 9 7]

```

=====
Iteration 1#:[4 6 3 2 1 9 7 ]
item moved : 6
item moved : 4
item inserted : 3, at position : 0
Iteration 2#:[3 4 6 2 1 9 7 ]
item moved : 6
item moved : 4
item moved : 3
item inserted : 2, at position : 0
Iteration 3#:[2 3 4 6 1 9 7 ]
item moved : 6
item moved : 4
item moved : 3
item moved : 2
item inserted : 1, at position : 0
Iteration 4#:[1 2 3 4 6 9 7 ]
Iteration 5#:[1 2 3 4 6 9 7 ]
item moved : 9
item inserted : 7, at position : 5
Iteration 6#:[1 2 3 4 6 7 9 ]
Output Array: [1 2 3 4 6 7 9 ]

```

Quick Sort: Quick sort is a highly efficient sorting algorithm and is based on partitioning of array of data into smaller arrays. A large array is partitioned into two arrays one of which holds values smaller than the specified value, say pivot, based on which the partition is made and another array holds values greater than the pivot value. Quick sort partitions an array and then calls itself recursively twice to sort the two resulting subarrays. This algorithm is quite efficient for large-sized data sets as its average and worst case complexity are of $O(n \log n)$, where n is the number of items.

Partition in Quick Sort:

Following example explains how to find the pivot value in an array.

The pivot value divides the list into two parts. And recursively, we find the pivot for each sub-lists until all lists contains only one element.

A[6]	A[7]	A[8]	A[9]	A[10]	A[11]	A[12]
98	84	65	108	60	96	72
72	84	65	108	60	96	98
72	84	65	98	60	96	108
72	84	65	96	60	98	108

Pivot

- Step 1 – Choose the highest index value as pivot
- Step 2 – Take two variables to point left and right of the list excluding pivot
- Step 3 – left points to the low index
- Step 4 – right points to the high index
- Step 5 – while value at left is less than pivot move right
- Step 6 – while value at right is greater than pivot move left
- Step 7 – if both step 5 and step 6 does not match swap left and right
- Step 8 – if $\text{left} \geq \text{right}$, the point where they met is new pivot

Quick Sort Algorithm:

Using pivot algorithm recursively, we end up with smaller possible partitions. Each partition is then processed for quick sort. We define recursive algorithm for quicksort as follows –

- Step 1 – Make the right-most index value pivot
- Step 2 – partition the array using pivot value
- Step 3 – quicksort left partition recursively
- Step 4 – quicksort right partition recursively

C Program for Quick Sort:

```
#include <stdio.h>
#include <stdlib.h>
#define MAX 7

int intArray[MAX] = {4,6,3,2,1,9,7};

void printline(int count) {
```

```

int i;

for(i = 0;i <count-1;i++) {
    printf("=");
}

printf("\n");
}

void display() {
    int i;
    printf("[");

    // navigate through all items
    for(i = 0;i<MAX;i++) {
        printf("%d ",intArray[i]);
    }

    printf("]\n");
}

void swap(int num1, int num2) {
    int temp = intArray[num1];
    intArray[num1] = intArray[num2];
    intArray[num2] = temp;
}

int partition(int left, int right, int pivot) {
    int leftPointer = left -1;
    int rightPointer = right;

    while(true) {
        while(intArray[++leftPointer] < pivot) {
            //do nothing
        }

        while(rightPointer > 0 && intArray[--rightPointer] > pivot) {
            //do nothing
        }

        if(leftPointer >= rightPointer) {
            break;
        }else {
            printf(" item swapped :%d,%d\n", intArray[leftPointer],intArray[rightPointer]);
            swap(leftPointer,rightPointer);
        }
    }
}

```

```

        printf(" pivot swapped :%d,%d\n", intArray[leftPointer],intArray[right]);
        swap(leftPointer,right);
        printf("Updated Array: ");
        display();
        return leftPointer;
    }

void quickSort(int left, int right) {
    if(right-left <= 0) {
        return;
    }else {
        int pivot = intArray[right];
        int partitionPoint = partition(left, right, pivot);
        quickSort(left,partitionPoint-1);
        quickSort(partitionPoint+1,right);
    }
}

main() {
    printf("Input Array: ");
    display();
    printline(50);
    quickSort(0,MAX-1);
    printf("Output Array: ");
    display();
    printline(50);
}

```

If we compile and run the above program, it will produce the following result –
Output

Input Array: [4 6 3 2 1 9 7]

```

=====
pivot swapped :9,7
Updated Array: [4 6 3 2 1 7 9 ]
pivot swapped :4,1
Updated Array: [1 6 3 2 4 7 9 ]
item swapped :6,2
pivot swapped :6,4
Updated Array: [1 2 3 4 6 7 9 ]
pivot swapped :3,3
Updated Array: [1 2 3 4 6 7 9 ]
Output Array: [1 2 3 4 6 7 9 ]

```

Important Points

- Sorting algorithm specifies the way to arrange data in a particular order. Most common orders are in numerical or lexicographical order.
- Bubble sort is a simple sorting algorithm. This sorting algorithm is comparison-based algorithm in which each pair of adjacent elements is compared and the elements are swapped if they are not in order.
- Merge sort is a sorting technique based on divide and conquer technique. With worst-case time complexity being $O(n \log n)$, it is one of the most respected algorithms.
- Quick sort is a highly efficient sorting algorithm and is based on partitioning of array of data into smaller arrays.

Exercise

Objective type questions.

- Q1. The complexity of Bubble sort algorithm is
- a. $O(n)$
 - b. $O(\log n)$
 - c. $O(n^2)$
 - d. $O(n \log n)$
- Q2. The complexity of merge sort algorithm is
- a. $O(n)$
 - b. $O(\log n)$
 - c. $O(n^2)$
 - d. $O(n \log n)$
- Q3. The complexity of selection sort algorithm is
- a. $O(n)$
 - b. $O(\log n)$
 - c. $O(n^2)$
 - d. $O(n \log n)$
- Q4. Which is the good sorting algorithm.
- a. Selection sort
 - b. Insertion sort
 - c. Quick sort
 - d. None
- Q5. The complexity of quick sort algorithm is.
- a. $O(n)$
 - b. $O(\log n)$
 - c. $O(n^2)$
 - d. $O(n \log n)$

Short answer type questions.

- Q1. What is sorting ?
- Q2. What is Stable sort ?
- Q3. What is in-place sorting ?
- Q4. What is worst case running time of Quick sort ?

Q5. When is the worst case for quick sort ?

Essay type questions.

Q1. Explain merge sort in detail ?

Q2. Which is the best sorting algorithm and Why ?

Q3. Explain Quick sort ?

Q4. Differentiate between selection and insertion sort ?

Q5. What is the difference between stable and unstable sorting ?

Answers

Ans1. C

Ans2. d

Ans3. c

Ans4. c

Ans5. d