Linean Programming

- Linear programming: Linear programming (LP) is an optimisation technique in which a linear function is optimisied (i.e. minimised or maximised) subject to certain constraints which are in the form of linear inequalities and equations. The function to be optimised is called objective function.
- Applications of linear programming: Linear programming optimum combination of several variables subject to certain constraints or restrictions.
- Formation of linear programming problem (LPP): The basic problem in the formulation of a linear programming problem is to set-up some mathematical model. This can be done by asking the following questions:
- (a) What are the unknown (variables)?
- (b) What is the objective?
- (c) What are the nestrictions?

FOH this, let $x_1, x_2, x_3, \dots, x_n$ be the vaniables. Let the objective function to be optimized (i.e. minimised on maximised) be given by Z.

- (i) $Z = C_1 x_1 + C_2 x_2 + \dots + C_n x_n$ where $c_i x_i$ ($i = 1, 2, \dots, n$) are constraints.
- (ii) Let there be mn constants and let a be a set of constants such that

$$a_{11} x_1 + a_{12} x_2 + \dots + a_n x_n (\leq , = on \geq) b_1$$
 $a_{21} x_1 + a_{22} x_2 + \dots + a_n x_n (\leq , = on \geq) b_2$
 \dots
 $a_{m1} x_1 + a_{m2} x_1 + \dots + a_{mn} x_n (\leq , = on \geq) b_m$

The problem of determinating values of x_1, x_2, \ldots, x_n which makes Z, a minimum or maximum and which satisfies (ii) and (iii) is called the general linear programming problem

- J General LPP
- (a) Decision variables: The variables $x_1, x_2, x_3, \dots, x_n$ whose values are to be decided, are called decision variables.
- (b) Objective function: The linear function $Z = c_1x_1 + c_2x_2 + \ldots + c_nx_n$ which is to be optimized (maximised on minimised) is called the objective function or preference function of the general linear programming problem.
- (c) Structural constraints: The inequalities given in (ii) are called the structural constraints of the general linear programming problem. The structural constraints are generally in the form of inequalities of ≥ type or ≤ type, but occasionally, a structural constraint may be in the form of an equation.
- (d) Non-negative constraints: The set of inequalities (iii) is usually known as the set of non-negative constraints: Constraints of the general LPP. These constraints imply that the variables x_1, x_2, \ldots, x_n cannot take negative values.
- (e) Feasible Solution: Any solution of a general LPP which satisfies all the constraints, structural and non negative, of the problem, is called a fesible solution of general LPP.
- (f) Optimum solution: Any feasible solution which optimizes (i.e. minimize on maximises) the objective function of the LPP is called optimum solution.

- Requirements for Mathematical Formulation of LPP: Before getting the mathematical form of a linear programming problem, it is important to recognize the problem which can be handled by linear programming problem. For the formulation of a linear programming problem, the problem must satisfy the following requirements:
- (i) There must be an objective to minimise or maximise something. The objective must be capable of being cleanly defined mathematically as a linear function.
- (ii) There must be alternative sounces of action so that the problem of selecting the best course of actions may anise.
- (iii) The nesounces must be in economically quantifiable limited supply. The gives the constnaints to LPP
- (iv) The constraints (nestrictions) must be capable of being expressed in the form of linear equations or inequalities.
- Solving linean Programming problem: To solve linear programming problems, Corner Point method is adopted. Under this method following steps are performed:
- Step 1: At finst, feasible negion is obtained by plotting the graph of given linear constraints and its connex points are obtained by solving the two equations of the lines intersecting at that point.
- Step II: The value of objective function Z = ax + by is obtained for each connen point by putling its x and y-coordinate in place of x and y in Z = ax + by. Let M and M be langest and smallest value of Z respectively.
 - Case I : If the feasible is bounded, then M and m are the maximum and minimum values of Z. Case II : If the feasible is unbounded, then we proceed as follows:
- Step III: The open half plane determined by ax + by > M and ax + by < m are obtained.
 - Case I: If there is no common point in the half plane determined by ax + by > M and feasible negion, then M is maximum value of Z, otherwise Z has no maximum value.
 - Case II: If there is no common point in the half plane determined by $ax + by \le M$ and feasible negion, then m is minimum value of Z, otherwise Z has no minimum value.